

Chapter 1

Welcome Aboard

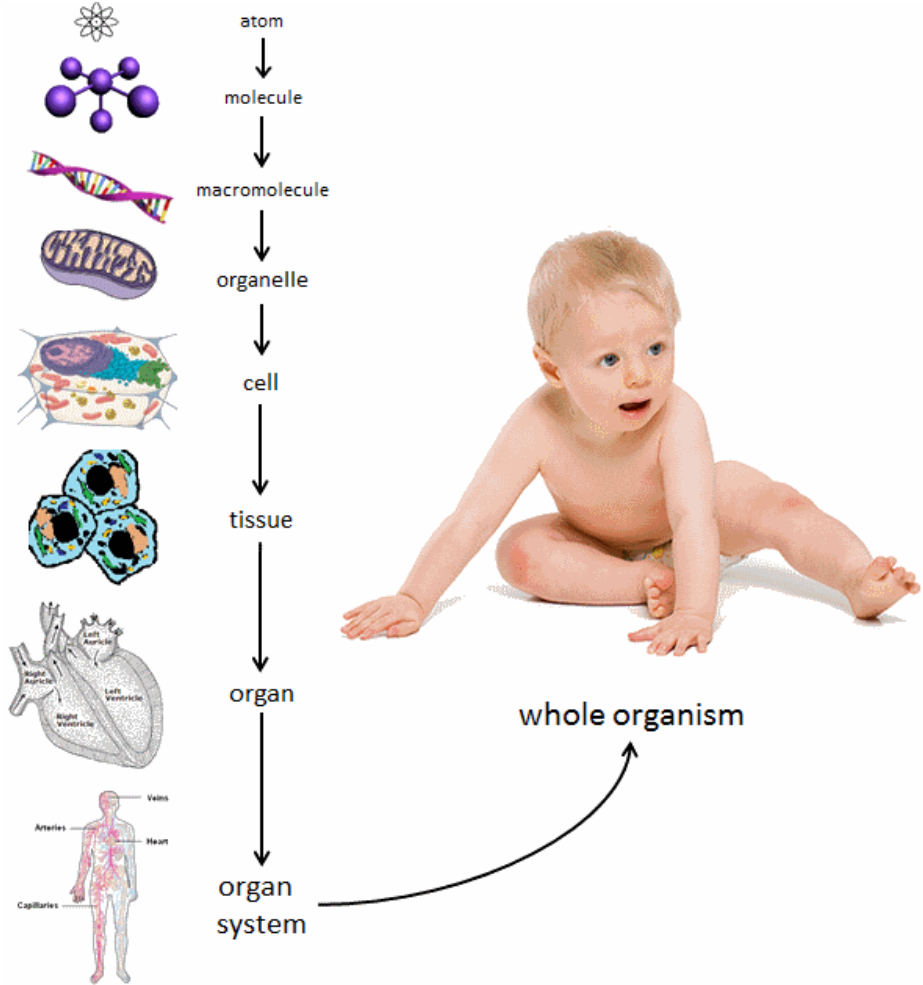
Abstraction



-----Interface



Levels of Abstraction (Biological System)



Source: https://media.studyisland.com/pics/40040_bio_organization.gif

Levels of Abstraction (Computer)

Problems

Algorithms

Language

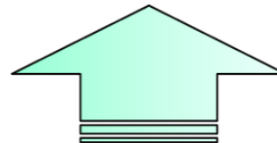
Instruction Set Architecture

Hardware/Software Interface

Microarchitecture

Circuits

Devices



Universal Computing Device

All computers, given enough time and memory, are capable of computing exactly the same things.



Embedded Processor

=



=



Supercomputer



Then what is the simplest possible computing device?

Turing Machine

Mathematical model of a device that can perform any computation – Alan Turing (1937)

Every computation can be performed by some Turing machine. (*Turing's thesis*)

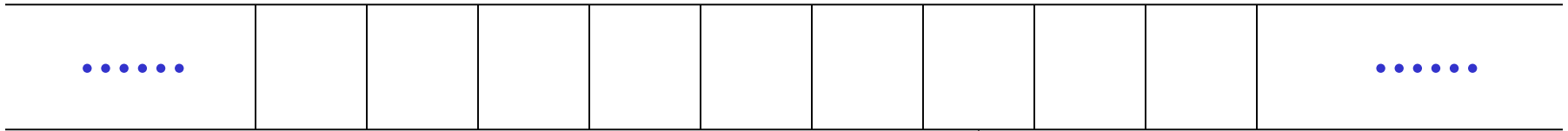
(그럼 이 세상에 계산하지 못하는 것도 있나? Yes...Halting Problem...We'll discuss it later)

For more info about Turing machines, see http://www.wikipedia.org/wiki/Turing_machine/

For more about Alan Turing, see <http://www.turing.org.uk/turing/>

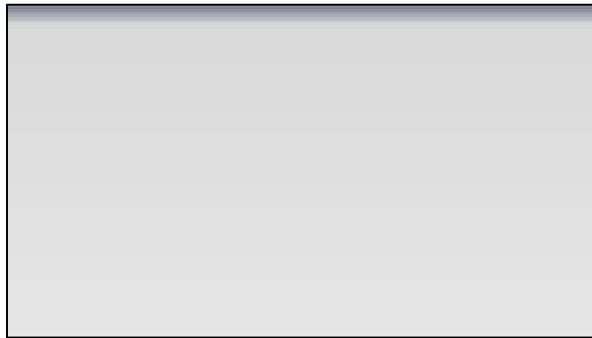
A Turing Machine

Tape



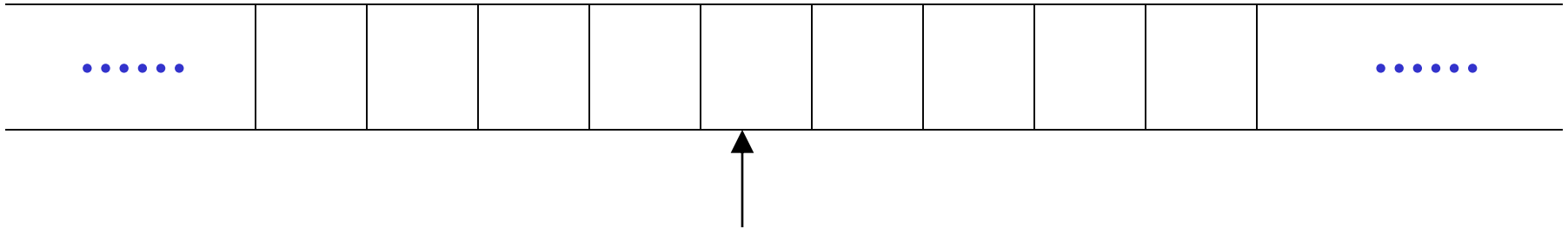
Read-Write head

Control Unit



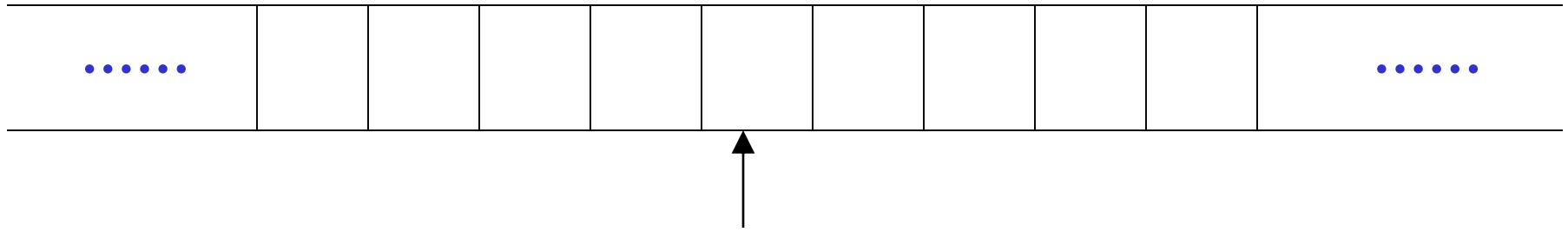
The Tape

No boundaries -- infinite length



Read-Write head

The head moves Left or Right



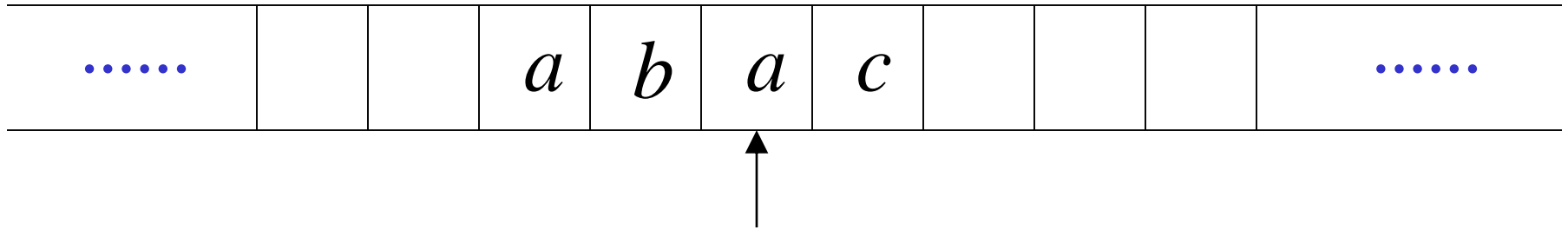
Read-Write head

The head at each transition (time step):

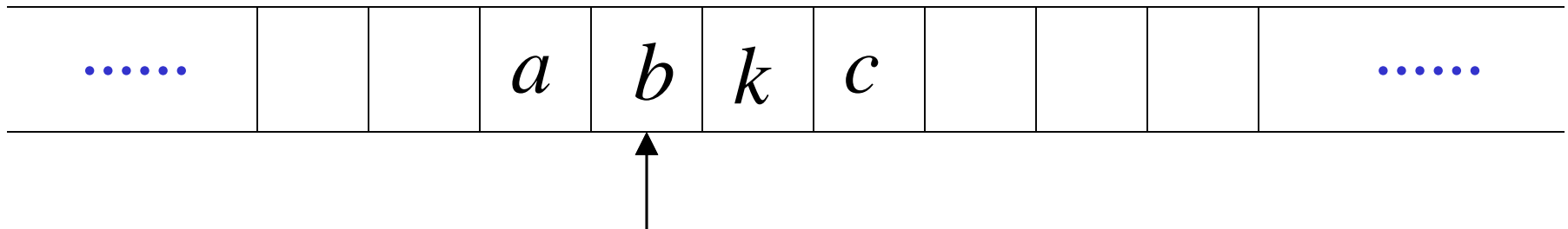
1. Reads a symbol
2. Writes a symbol
3. Moves Left or Right

Example:

Time 0



Time 1



1. Reads a

2. Writes k

3. Moves Left

Computation Example

The function $f(x, y) = x + y$ is computable

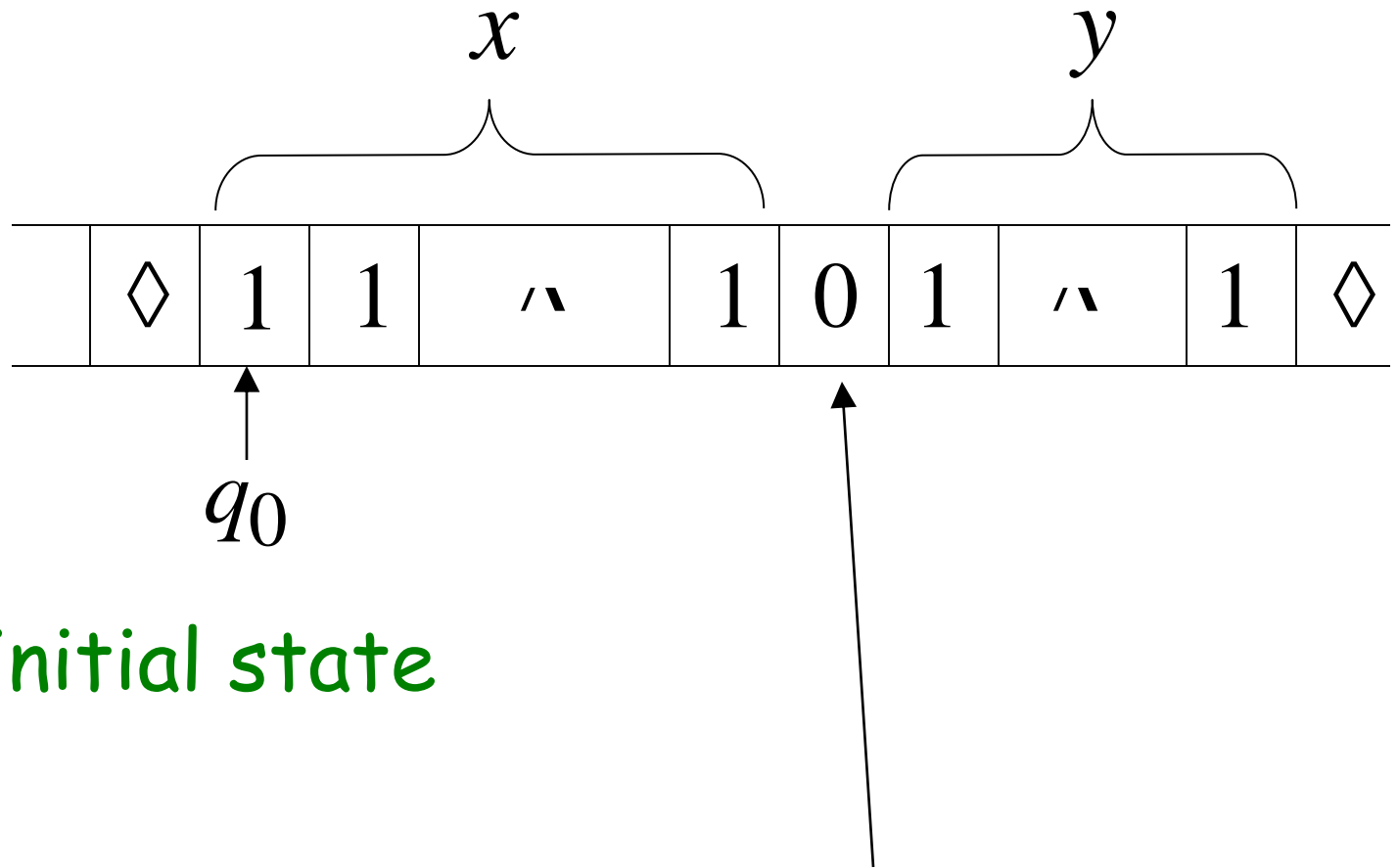
x, y are integers

Turing Machine:

Input string: $x0y$ unary

Output string: $xy0$ unary

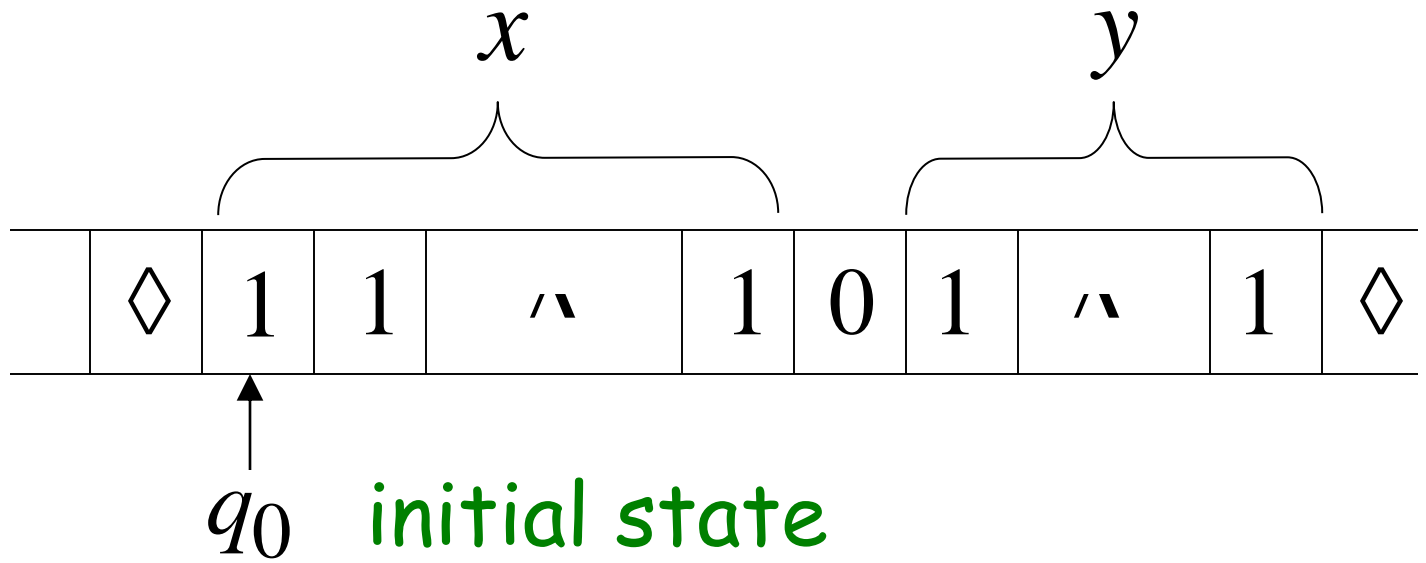
Start



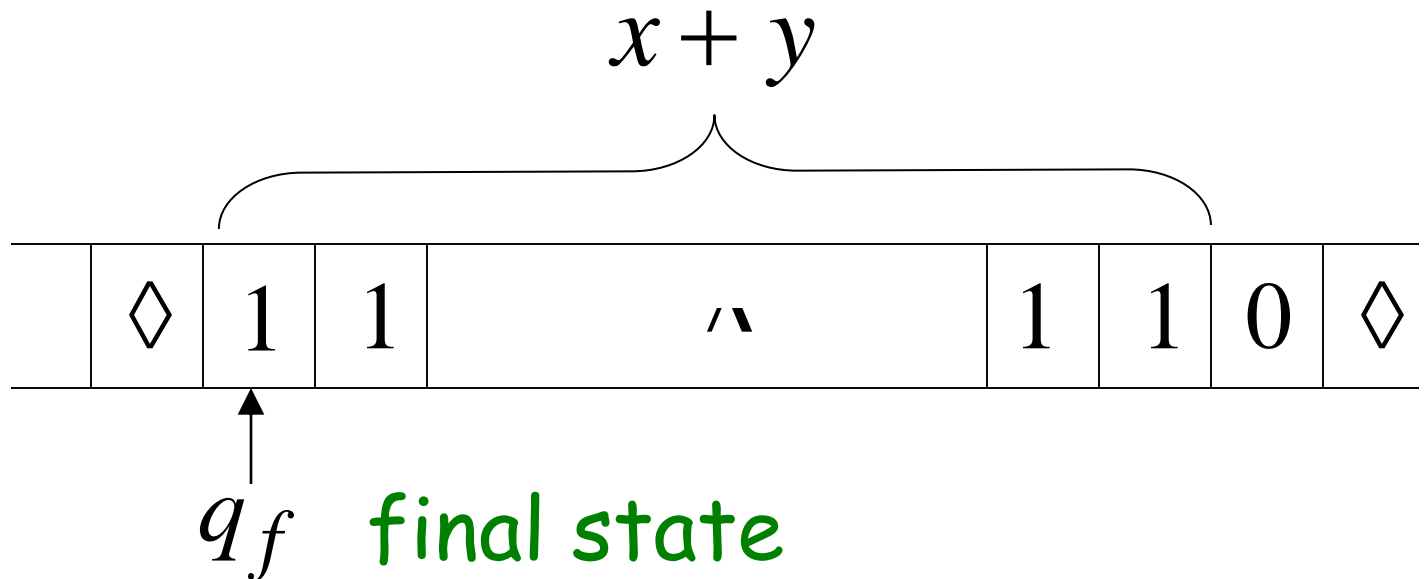
initial state

The 0 is the delimiter that separates the two numbers

Start



Finish

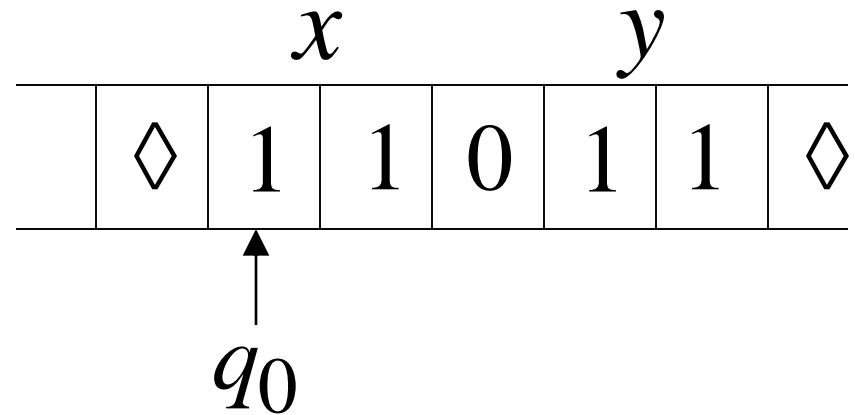


Execution Example:

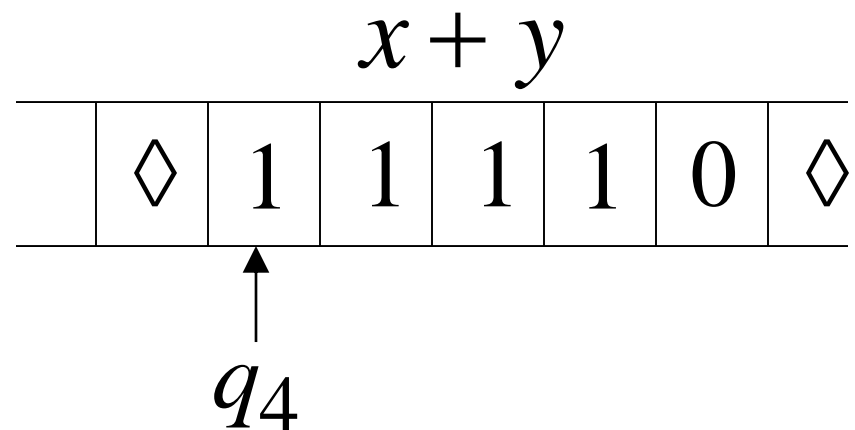
$$x = 11 \quad (=2)$$

$$y = 11 \quad (=2)$$

Time 0

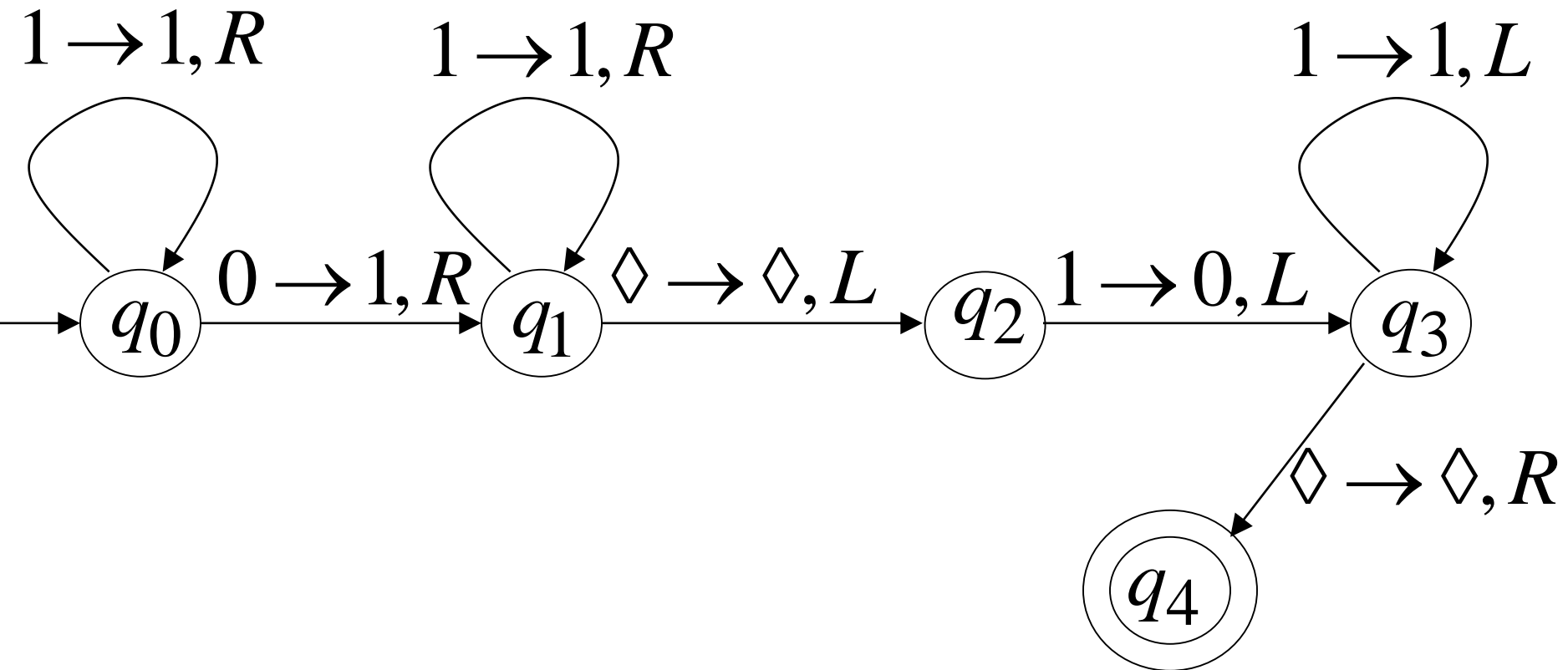


Final Result

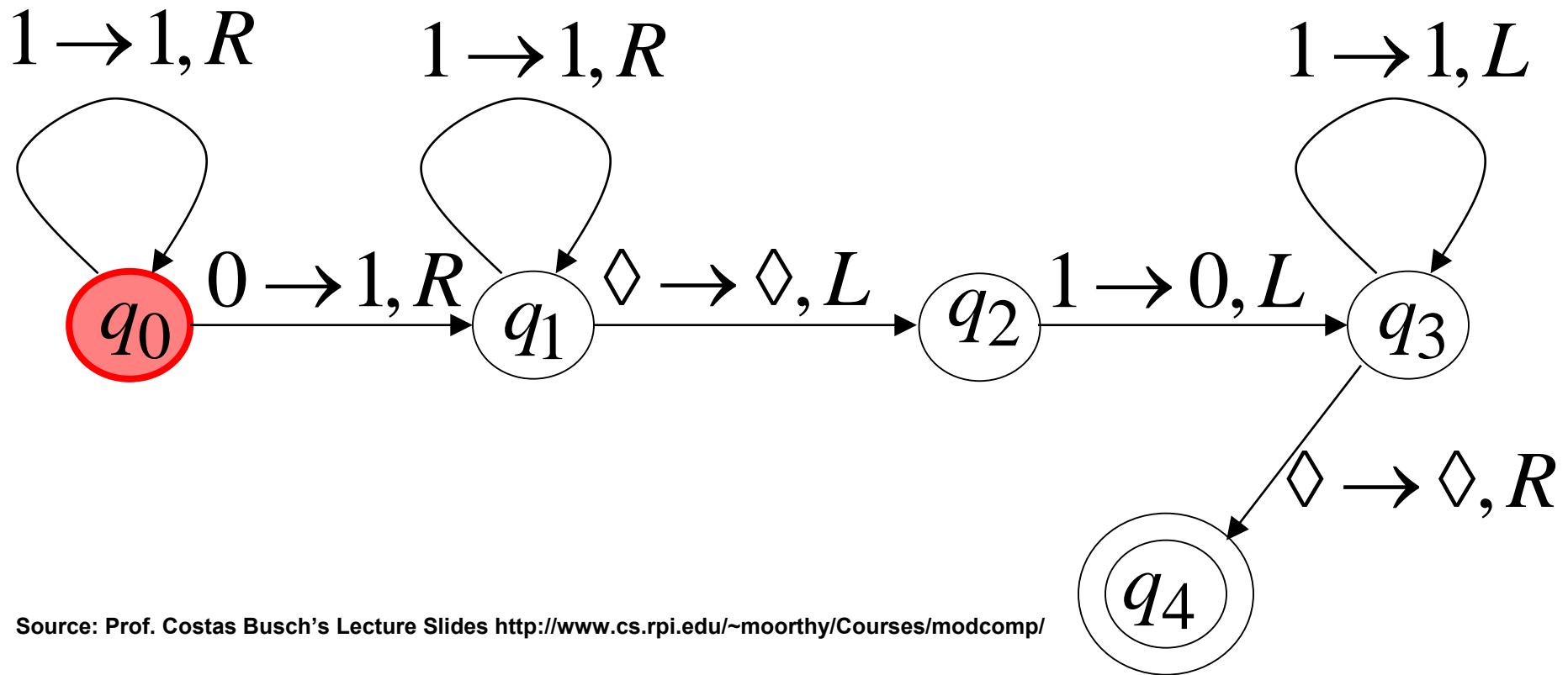
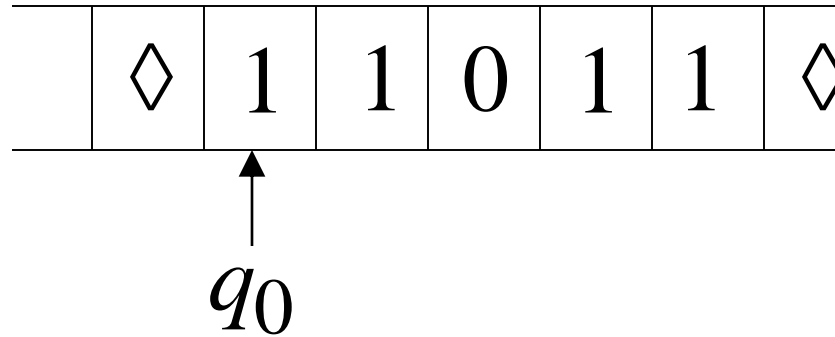


Turing machine for function $f(x, y) = x + y$

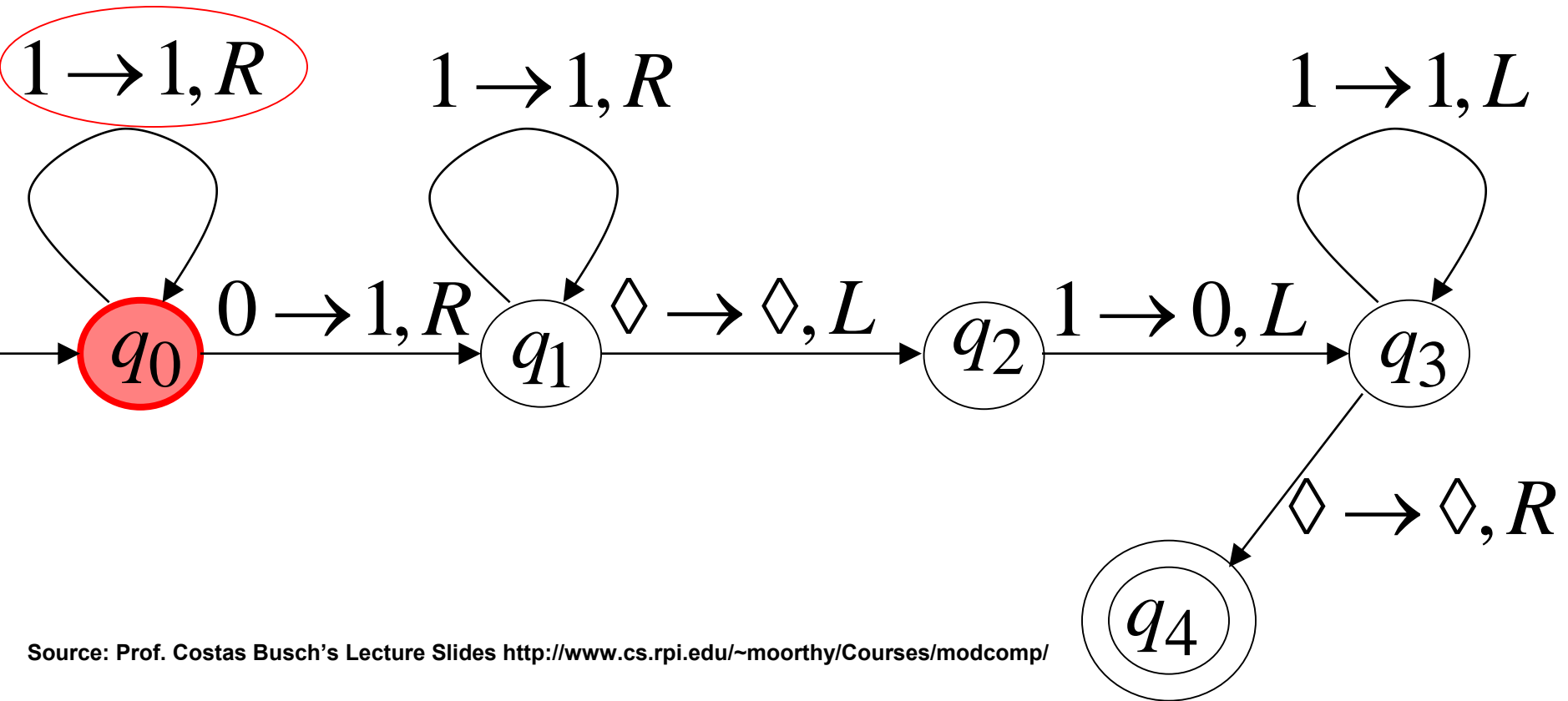
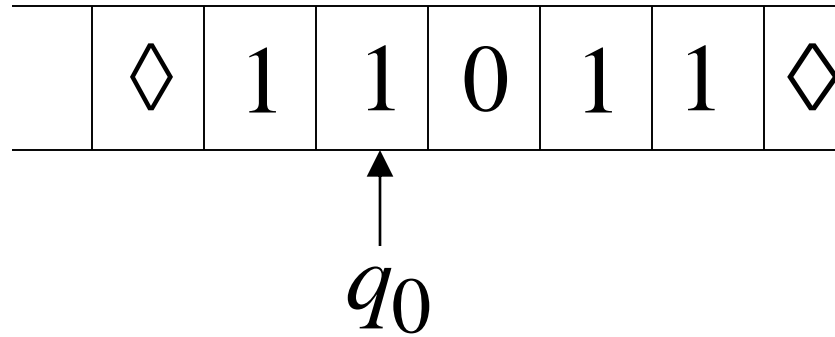
Control Unit



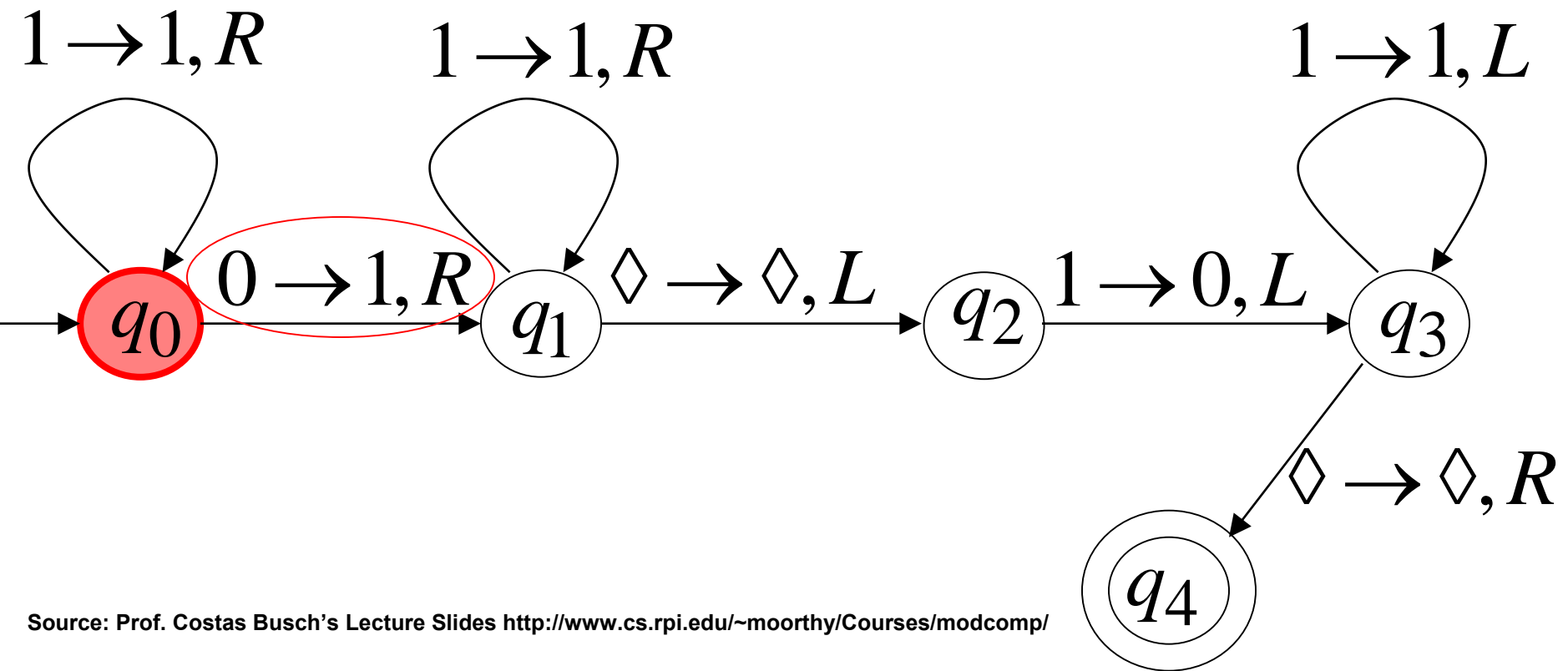
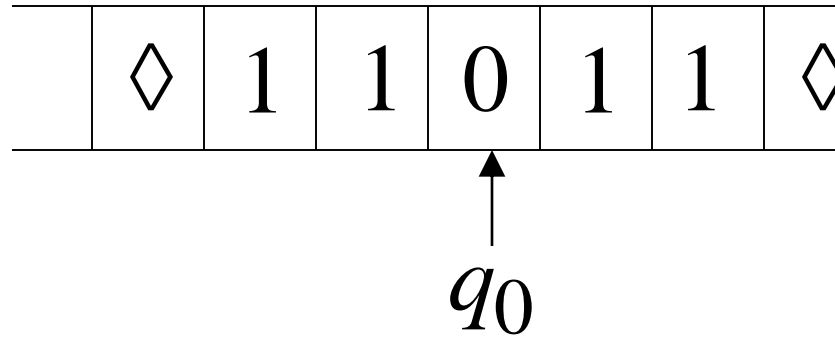
Time 0



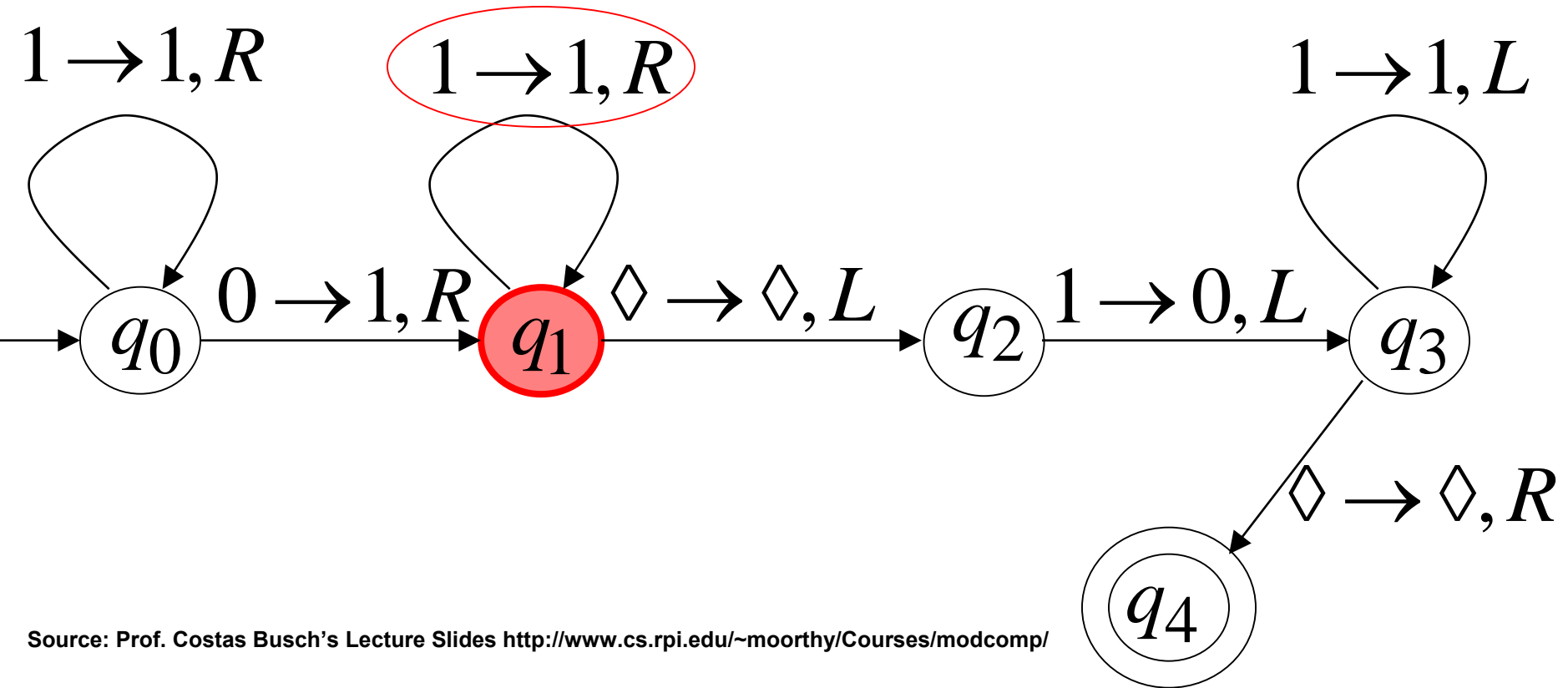
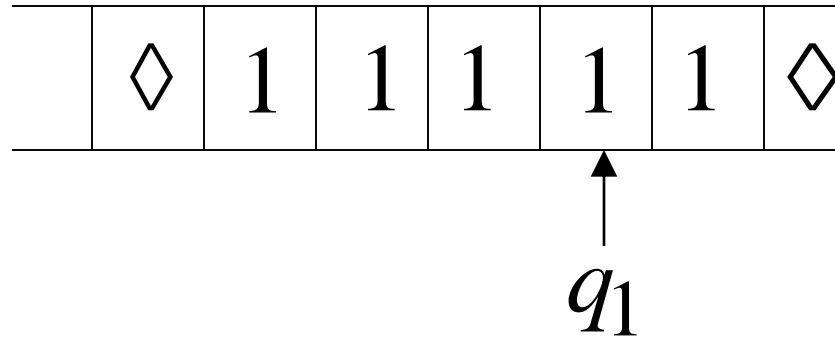
Time 1



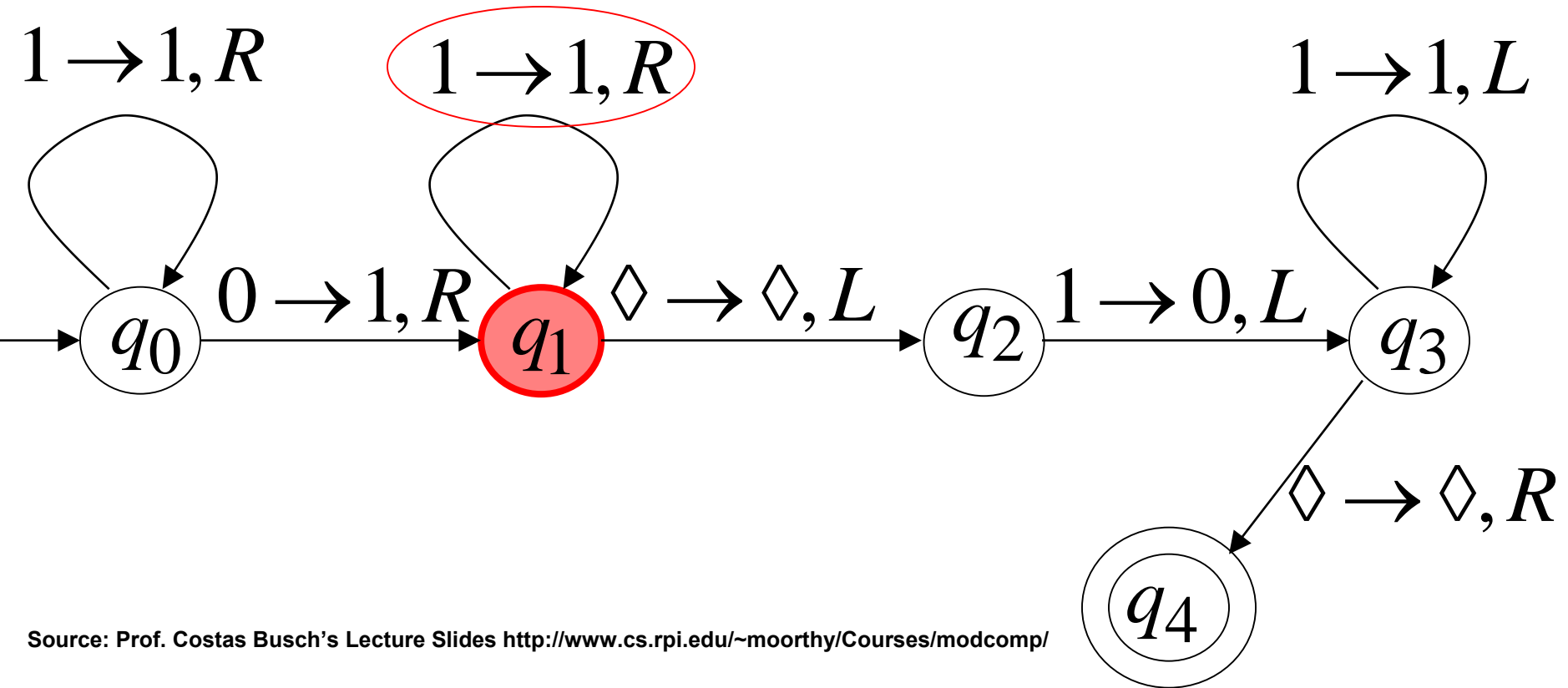
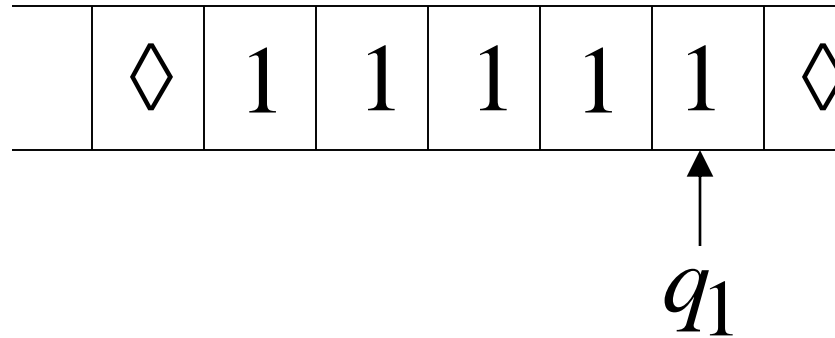
Time 2



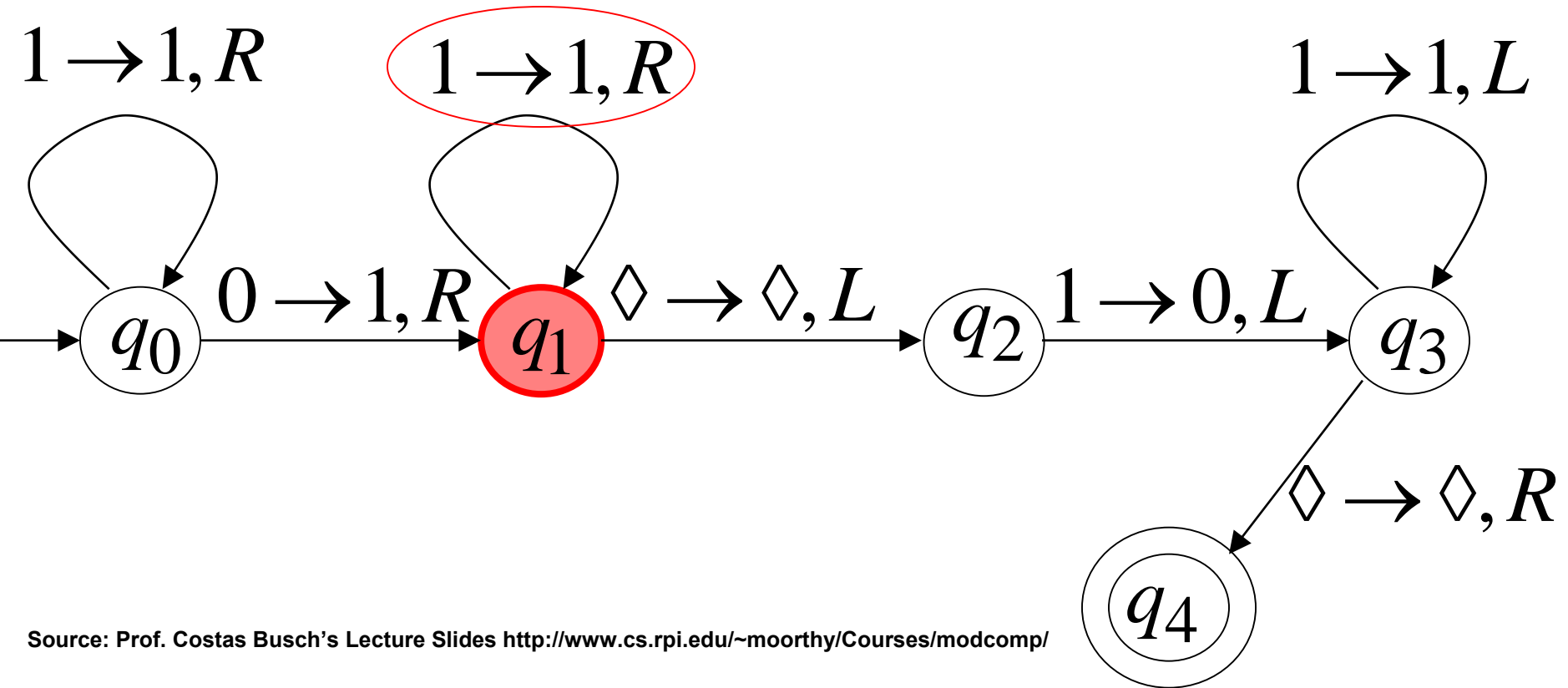
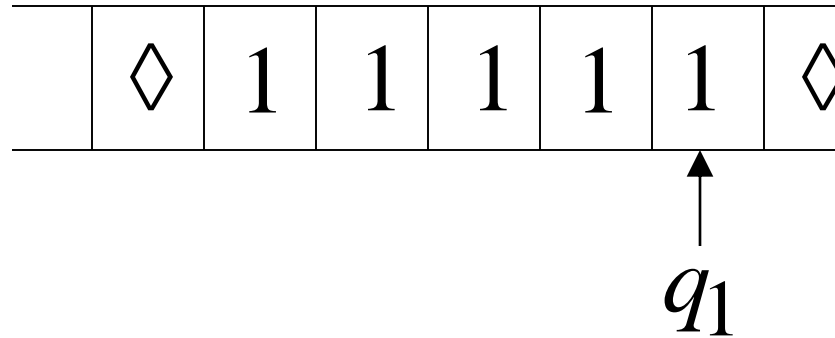
Time 3



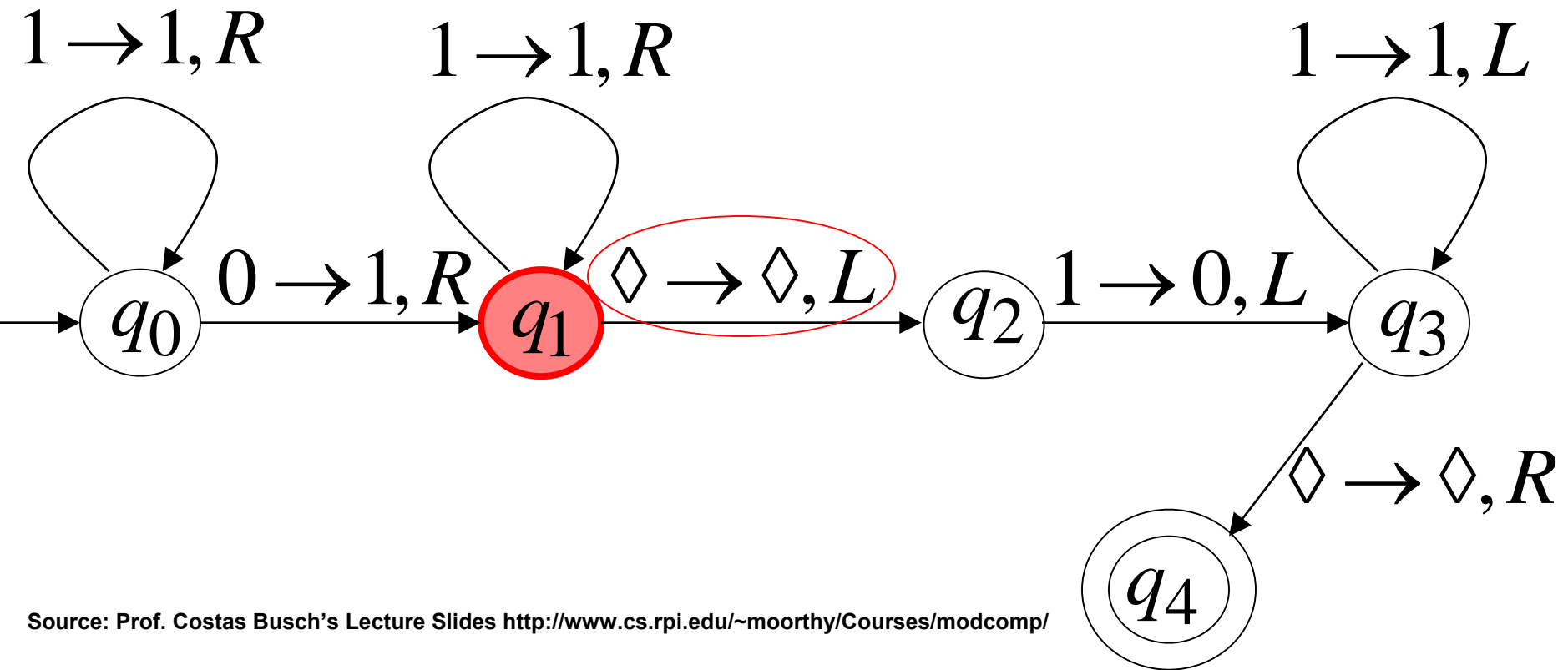
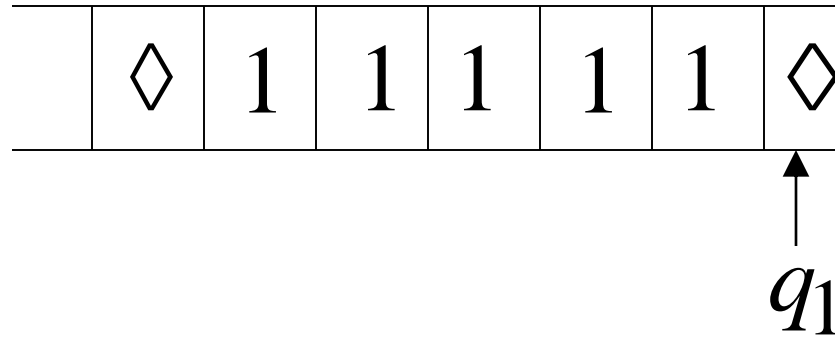
Time 4



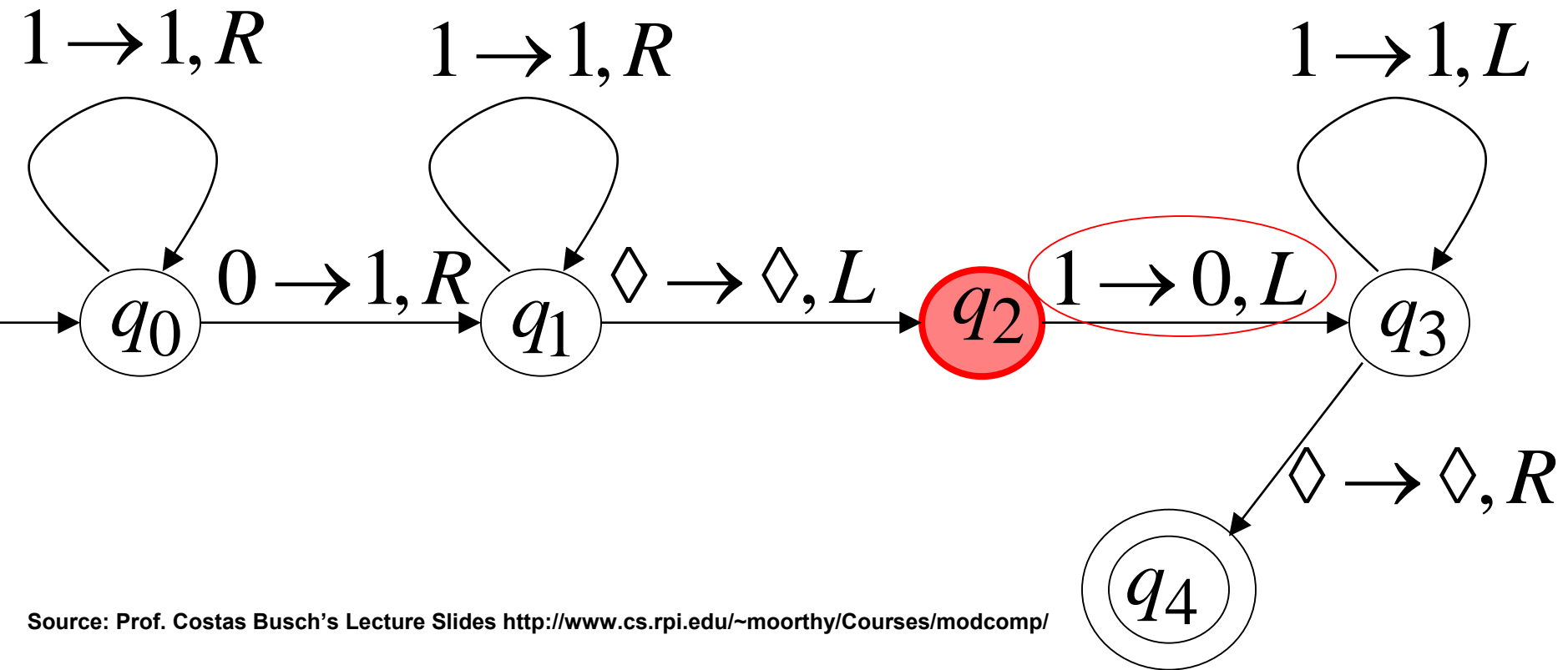
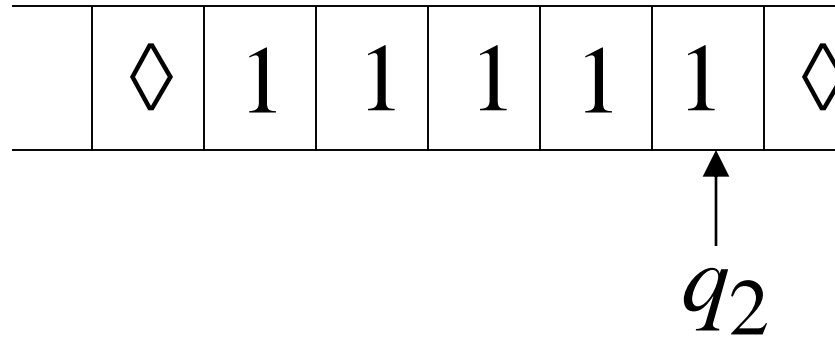
Time 4



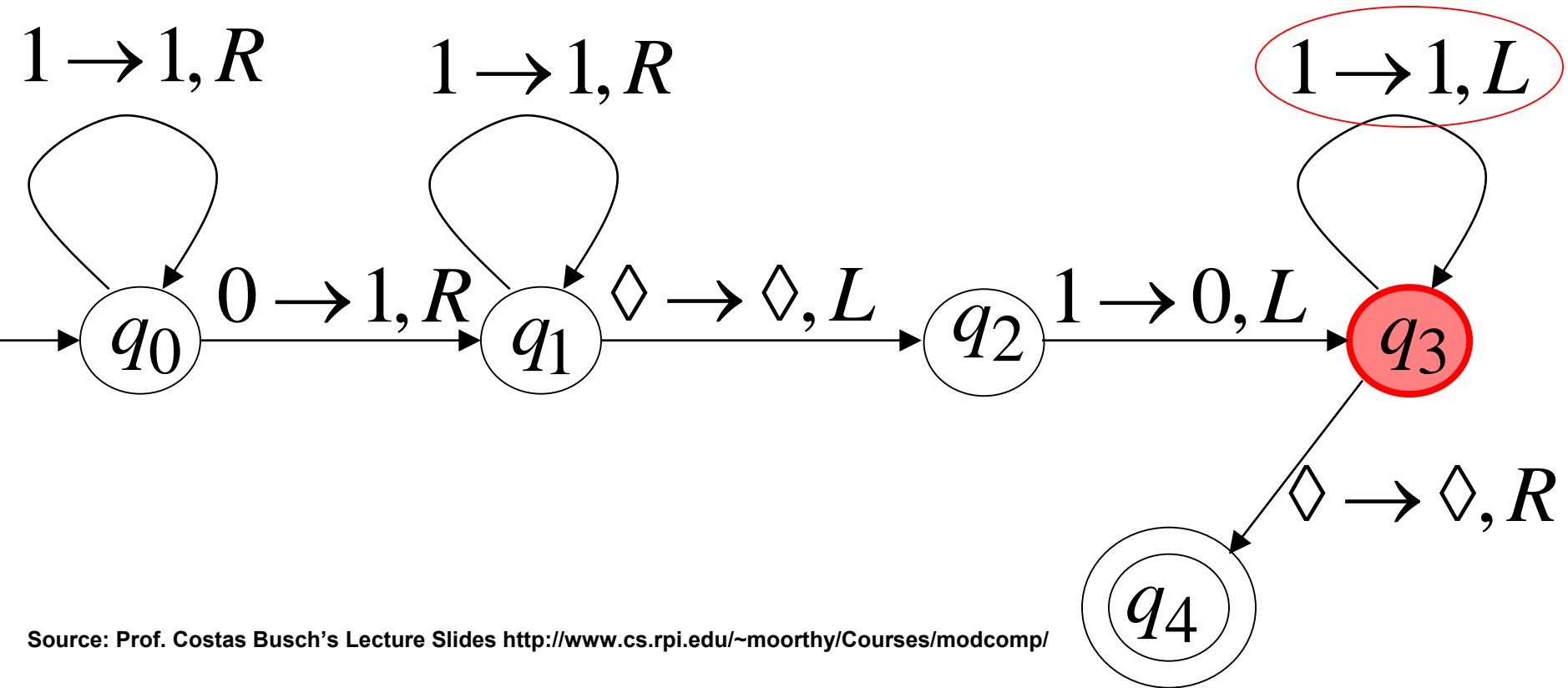
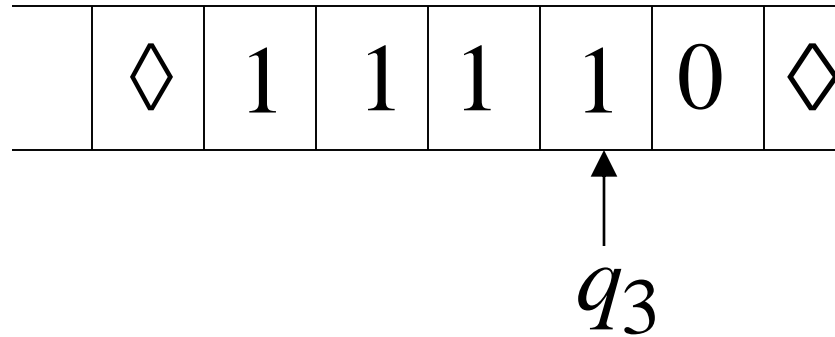
Time 5



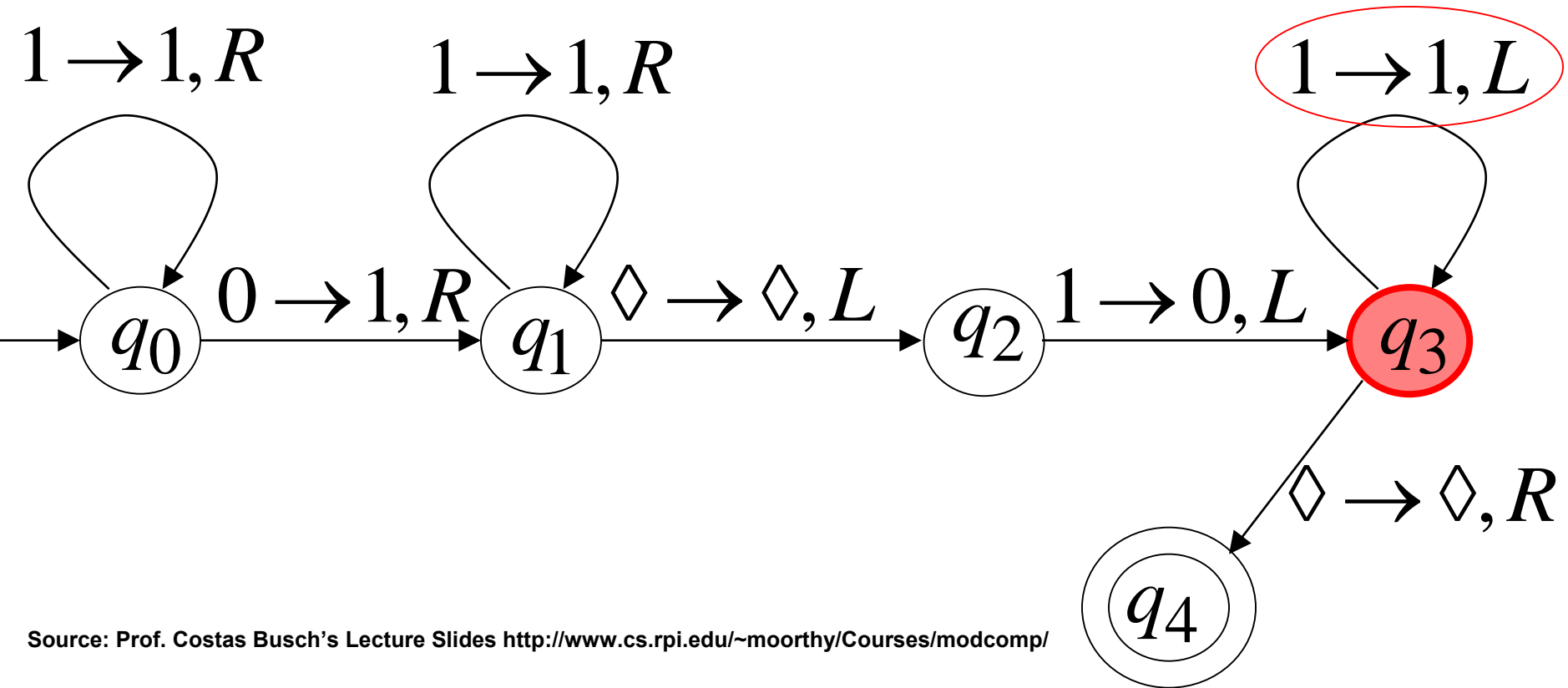
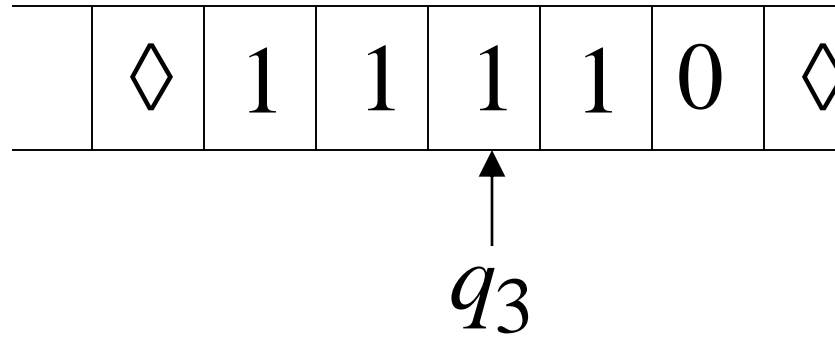
Time 6



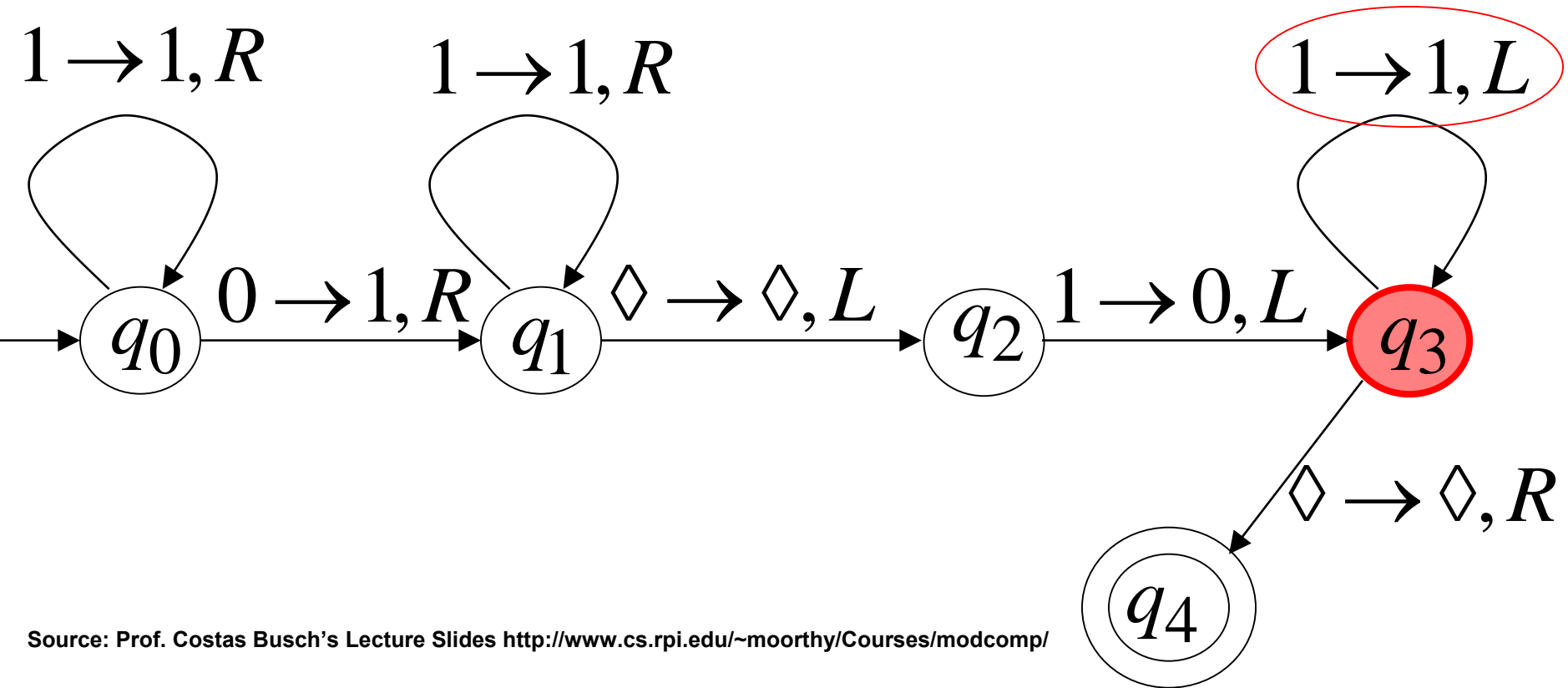
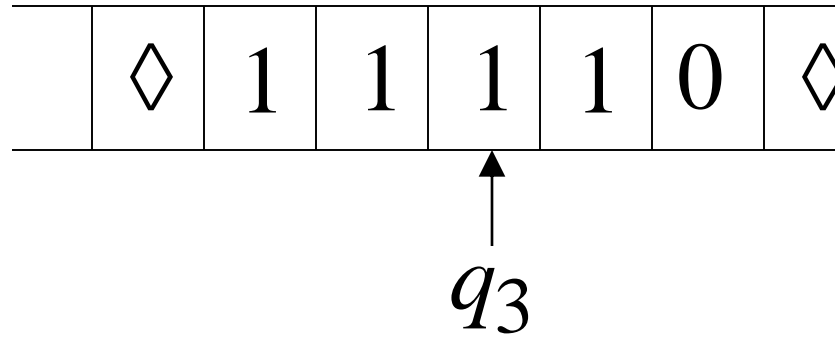
Time 7



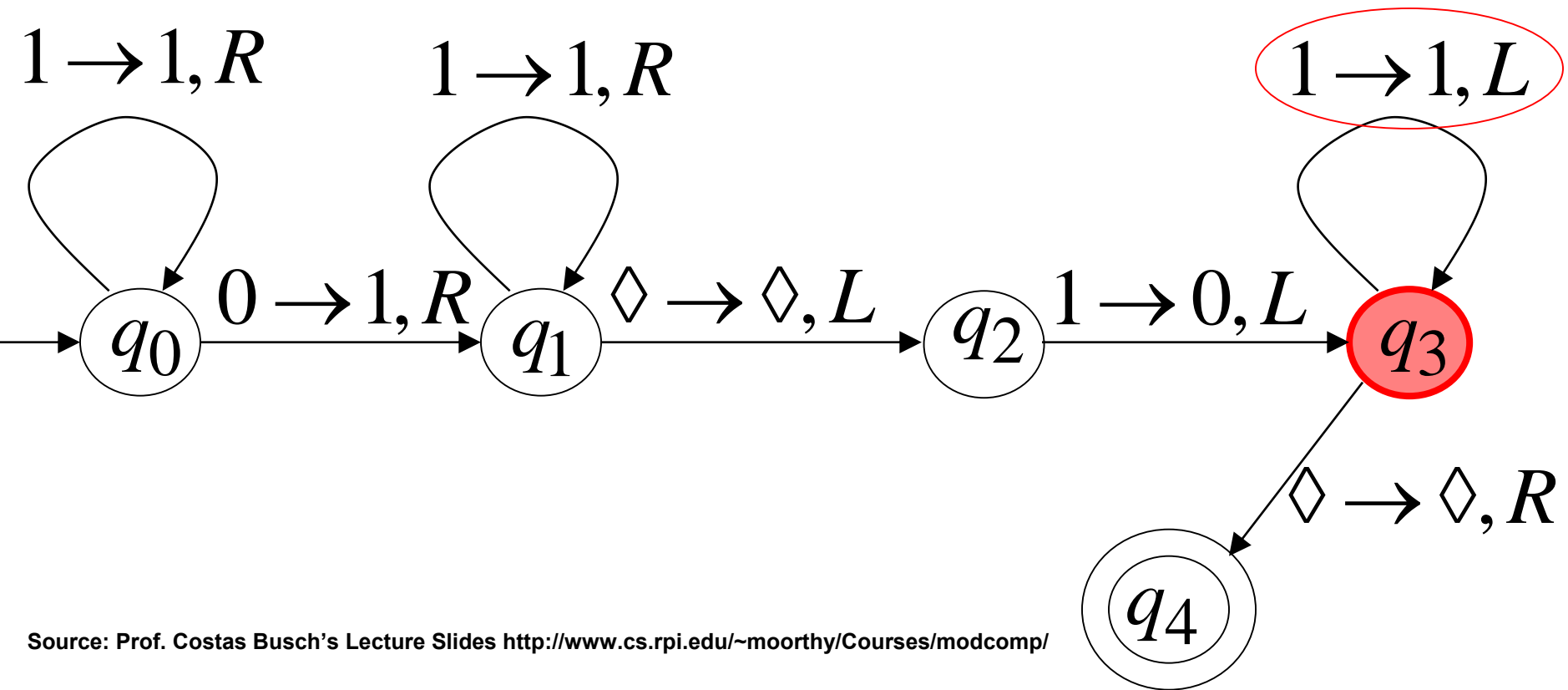
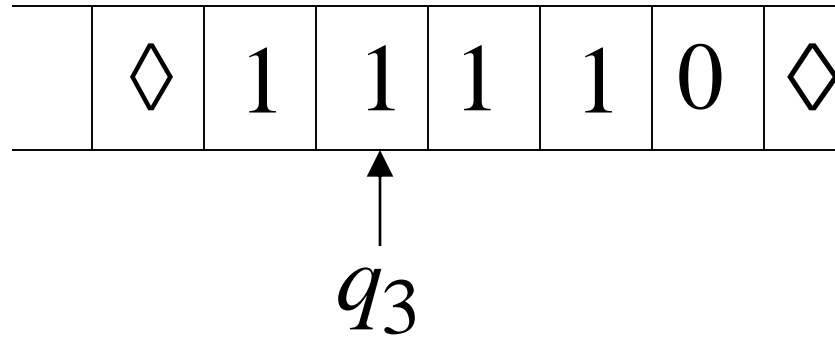
Time 8



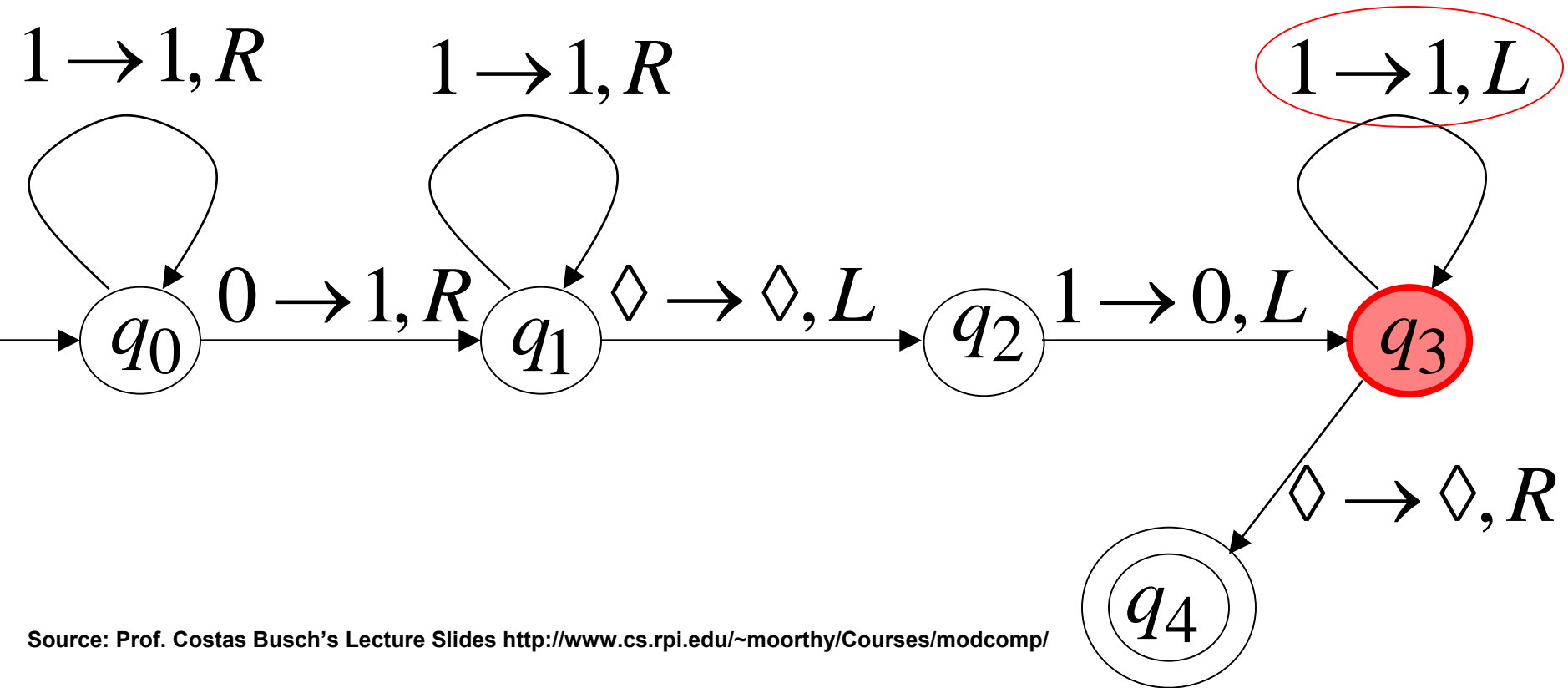
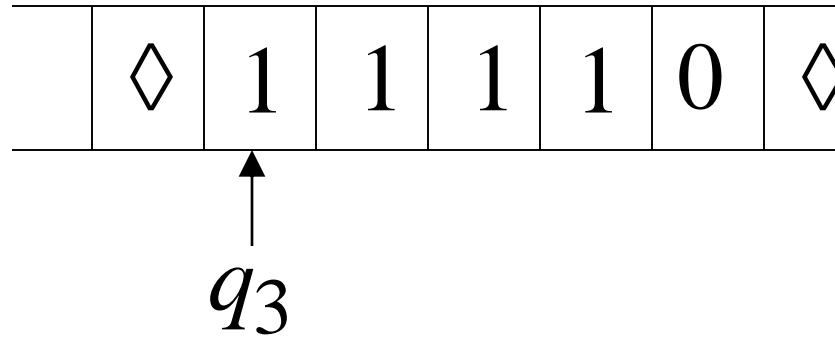
Time 8



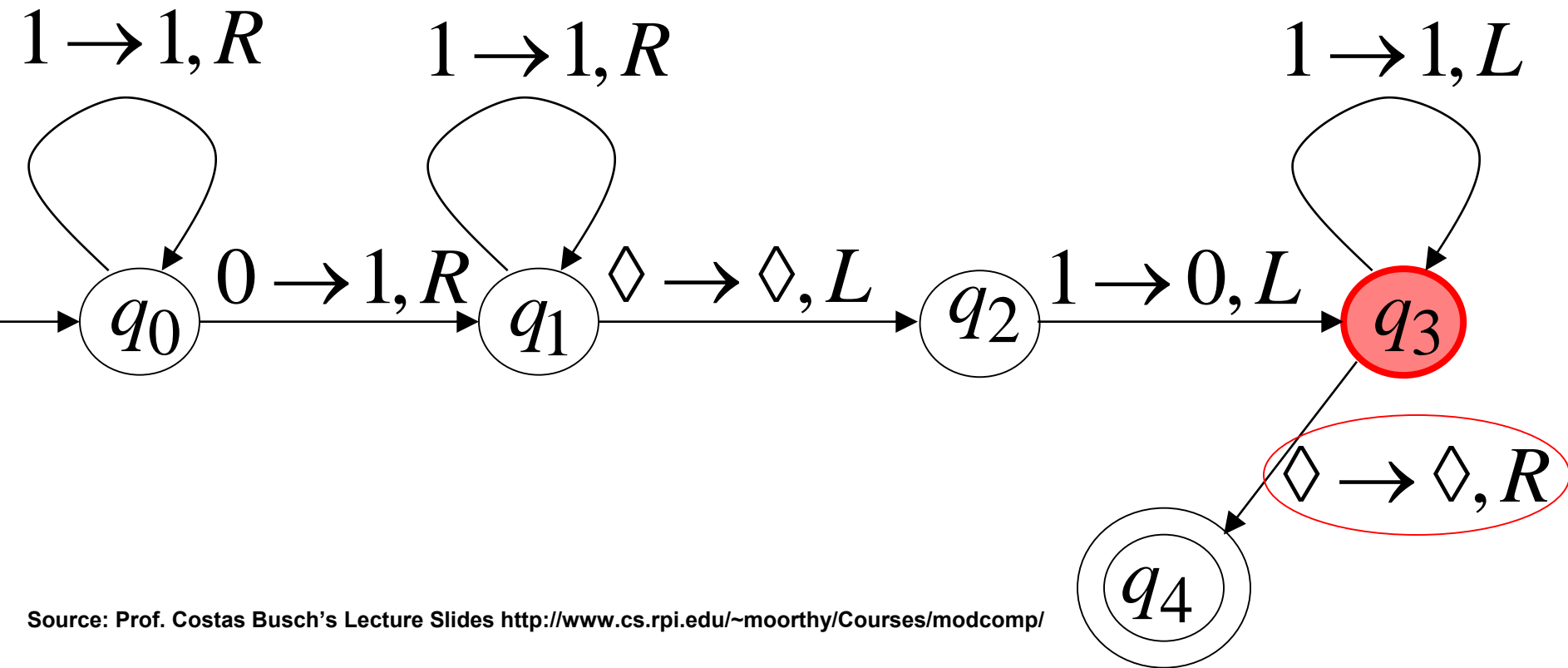
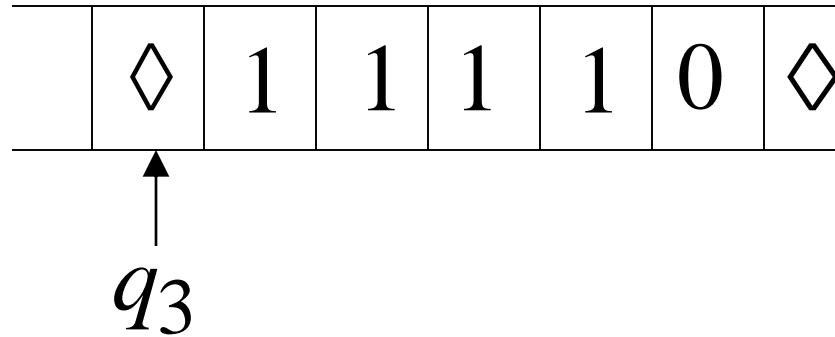
Time 9



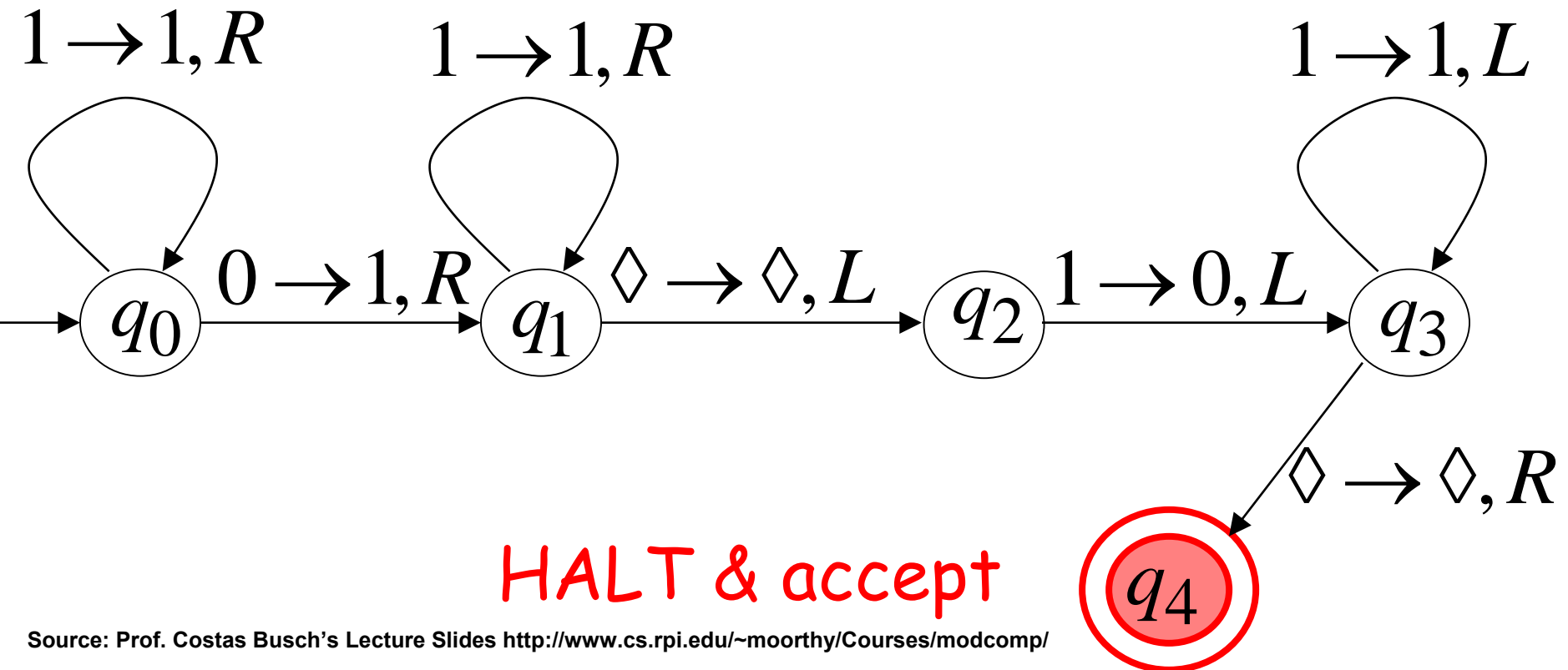
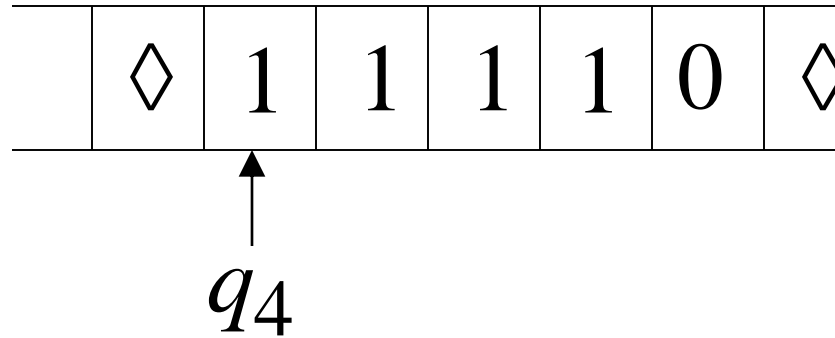
Time 10



Time 11



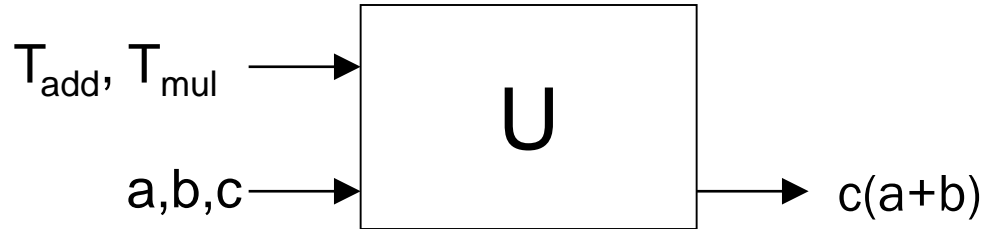
Time 12



Universal Turing Machine

A machine that can implement all Turing machines
-- this is also a Turing machine!

- inputs: data, plus a description of computation (other TMs)



Universal Turing Machine

U is programmable – so is a computer!

- Program is part of the input data
- a computer can emulate a Universal Turing Machine and vice versa

A computer is a universal computing device.

Halting Problem

- **Halting Problem**
 - **The problem of determining, from a description of an arbitrary computer program (i.e., Turing machine) and an input, whether the program will finish running (i.e., halts) or continue to run forever**
- **Halting problem is undecidable (not Turing machine solvable)**
(Proof by an application of Cantor's diagonal argument)
오토마타 교과목에서 더 깊게 다루어짐.

꼭 기억해야 할 것

- (Levels of) Abstraction
- Turing Equivalence
- Undecidable Problem (not Turing machine computable)